

# A NOTE ON A DEGREE SUM CONDITION FOR LONG CYCLES IN GRAPHS

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ABSTRACT. We conjecture that a 2-connected graph  $G$  of order  $n$ , in which  $d(x) + d(y) \geq n - k$  for every pair of non-adjacent vertices  $x$  and  $y$ , contains a cycle of length  $n - k$  ( $k < n/2$ ), unless  $G$  is bipartite and  $n - k$  is odd. This generalizes to long cycles a well-known degree sum condition for hamiltonicity of Ore. The conjecture is shown to hold for  $k = 1$ .

## 1. INTRODUCTION

The subject of this note is the following conjecture, in which we generalize to long cycles a well-known degree sum condition for hamiltonicity of Ore [4]. All graphs considered are finite, undirected, with no loops or multiple edges.

**Conjecture 1.1.** *Let  $G$  be a 2-connected graph of order  $n \geq 3$ ,  $n \neq 5, 7$ , and let  $k < n/2$  be an integer. If*

$$d(x) + d(y) \geq n - k$$

*for every pair of non-adjacent vertices  $x$  and  $y$ , then  $G$  contains a cycle of length  $n - k$ , unless  $G$  is bipartite and  $n - k \equiv 1 \pmod{2}$ .*

**Remark 1.2.** The conjecture is sharp. First of all, a quick look at  $C_5$  and  $C_7$  ensures that the assumption  $|G| \neq 5, 7$  is necessary. Secondly, it is easy to see that without the 2-connectedness assumption, there could be no long cycles at all. Consider, for instance, a graph  $G$  obtained from disjoint cliques  $H_1 = K_{\lfloor n/2 \rfloor}$  and  $H_2 = K_{\lceil n/2 \rceil}$  by joining a single vertex  $x_0$  of  $H_2$  with every vertex of  $H_1$ . Finally, the bound for the degree sum of non-adjacent vertices is best possible, as shown in the example below.

**Example 1.3.** Let  $G$  be a graph obtained from the complete bipartite graph  $K_{(n-k-1)/2, (n+k+1)/2}$  by joining all the vertices in the smaller colour class. Then  $d(x) + d(y) \geq n - k - 1$  for every pair of non-adjacent vertices  $x$  and  $y$ , and  $G$  contains no cycle of length greater than  $n - k - 1$ .

Our main result is the following theorem that implies Conjecture 1.1 for  $k = 1$ , as shown in Section 2. The proof of Theorem 1.4 is given in the last section.

**Theorem 1.4.** *Let  $G$  be a 2-connected graph of order  $n \geq 3$ , in which*

$$d(x) + d(y) \geq n - 1$$

*for every pair of non-adjacent vertices  $x$  and  $y$ .*

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- (i) If  $n$  is even, then  $G$  is hamiltonian.
- (ii) If  $n$  is odd, then  $G$  contains a cycle of length at least  $n - 1$ .

Moreover,  $G$  is not hamiltonian only if the minimal degree of its  $n$ -closure,  $Cl_n(G)$ , equals  $(n - 1)/2$ . In this case,  $Cl_n(G)$  is a maximal non-hamiltonian graph.

Recall that the  $n$ -closure  $Cl_n(G)$  of  $G$  is a graph obtained from  $G$  by successively joining all pairs  $(x, y)$  of non-adjacent vertices satisfying  $d(x) + d(y) \geq n$ .

## 2. LONG CYCLES IN GRAPHS

**Proposition 2.1.** *Conjecture 1.1 holds for  $k = 1$ .*

For the proof, we will need the following result of [3]:

**Theorem 2.2** (Haggkvist-Faudree-Schelp). *Let  $G$  be a hamiltonian graph on  $n$  vertices. If  $G$  contains more than  $\lfloor \frac{(n-1)^2}{4} \rfloor + 1$  edges, then  $G$  is pancyclic or bipartite.*

*Proof of Proposition 2.1.* By Theorem 1.4, we may assume that  $G$  is hamiltonian. Suppose first that  $G$  is a 2-connected non-bipartite hamiltonian graph of order  $n$ , in which  $d(x) + d(y) \geq n - 1$  whenever  $xy \notin E(G)$ .

Consider a vertex  $x$  of minimal degree  $d(x) = \delta(G)$  in  $G$ . Write  $\delta = \delta(G)$ . Then  $G$  has precisely  $n - 1 - \delta$  vertices non-adjacent to  $x$ , each of degree at least  $n - 1 - \delta$ . The remaining  $\delta + 1$  vertices are of degree at least  $\delta$  each, hence

$$\|G\| \geq \frac{1}{2}[(\delta + 1)\delta + (n - 1 - \delta)^2].$$

As  $\delta \geq 2$ , one immediately verifies that

$$\frac{1}{2}[(\delta + 1)\delta + (n - 1 - \delta)^2] > \frac{(n - 1)^2}{4} + 1,$$

whenever  $n \neq 5$ .

It remains to consider the case of  $G$  a bipartite 2-connected hamiltonian graph of order  $n$ . But then  $n$  must be even, for otherwise  $G$  would contain an odd cycle. Thus  $n - 1 \equiv 1 \pmod{2}$ , which completes the proof.  $\square$

For convenience, let us finally recall two well-known results, that we shall need in the proof of Theorem 1.4:

**Theorem 2.3** (Dirac [2]). *Let  $G$  be a graph of order  $n \geq 3$  and minimal degree  $\delta(G) \geq n/2$ . Then  $G$  is hamiltonian.*

**Theorem 2.4** (Bondy-Chvatal [1]). *Let  $G$  be a graph of order  $n$  and suppose that there is a pair of non-adjacent vertices  $x$  and  $y$  of  $G$  such that  $d(x) + d(y) \geq n$ . Then  $G$  is hamiltonian if and only if  $G + xy$  is hamiltonian.*

**Corollary 2.5.** *A graph  $G$  is hamiltonian if and only if its  $n$ -closure  $Cl_n(G)$  is so.*

## 3. PROOF OF THEOREM 1.4

*Proof of part (i).* Suppose there exists an even integer  $n \geq 4$  for which the assertion of the theorem does not hold. Let  $G$  be a maximal non-hamiltonian 2-connected graph of order  $n$ , in which  $d(x) + d(y) \geq n - 1$  whenever  $xy \notin E(G)$ .

By maximality of  $G$ ,  $G + xy$  is hamiltonian for every pair of non-adjacent vertices  $x, y \in V(G)$ . Hence, by Theorem 2.4, we must have

$$(*) \quad d(x) + d(y) = n - 1 \quad \text{whenever } xy \notin E(G).$$

The minimal degree  $\delta(G)$  of  $G$  satisfies inequality  $\delta(G) < n/2$ , by Theorem 2.3, hence, in particular,  $n - 1 - \delta(G) \geq \delta(G) + 1$ .

Pick  $x \in V(G)$  with  $d(x) = \delta(G)$ . There are precisely  $n - 1 - \delta(G)$  vertices in  $G$  non-adjacent to  $x$ , each of degree  $n - 1 - \delta(G)$ , by (\*). Put  $V = \{v \in V(G) : xv \notin E(G)\}$ . Pick  $y \in V$ . As  $d(y) = n - 1 - \delta(G)$ , there are precisely  $\delta(G)$  vertices in  $G$  non-adjacent to  $y$ , each of degree  $\delta(G)$ , by (\*) again. Put  $U = \{u \in V(G) : uy \notin E(G)\}$ . Then  $|U| = \delta(G)$ ,  $|V| = n - 1 - \delta(G)$ , and  $U \cap V = \emptyset$ , because vertices in  $U$  are of degree  $\delta(G)$  and those in  $V$  are of degree  $n - 1 - \delta(G) > \delta(G)$ . It follows that there exists a vertex  $z$  in  $G$  such that  $V(G) = U \cup V \cup \{z\}$  is a partition of the vertex set of  $G$ .

We will now show that  $d(z) = n - 1$ : Observe first that  $d(z) > \delta(G)$ . Indeed, if  $d(z) = \delta(G)$ , then by (\*),  $z$  is adjacent to every vertex in  $U$ , as  $2\delta(G) < n - 1$ . But  $z$  is also adjacent to  $y$ , as  $z \notin U$ , hence  $d(z) \geq |U| + 1 = \delta(G) + 1$ ; a contradiction. Consequently,  $z$  is adjacent to every vertex in  $V$ , by (\*) again, as  $d(z) + (n - 1 - \delta(G)) > n - 1$ . Hence  $d(z) \geq |V| = n - 1 - \delta(G)$ . On the other hand,  $z$  is adjacent to  $x$ , as  $z \notin V$ , which yields  $d(z) \geq |V| + 1 = n - \delta(G)$ . This last inequality paired with (\*) implies that  $z$  is adjacent to every other vertex in  $G$ , as required.

Next observe that  $u_1u_2 \in E(G)$  for every pair of vertices  $u_1, u_2$  in  $U$ , as  $d(u_1) + d(u_2) = 2\delta(G) < n - 1$ . It follows that  $N(u) \supset U \cup \{z\} \setminus \{u\}$ , and hence, by comparing cardinalities,  $N(u) = U \cup \{z\} \setminus \{u\}$  for every  $u \in U$ .

Similarly,  $v_1v_2 \in E(G)$  for every pair  $v_1, v_2$  in  $V$ , hence  $N(v) = V \cup \{z\} \setminus \{v\}$  for every  $v \in V$ . Therefore  $G = G_1 \cup G_2$ , where  $G_1$  is a complete graph of order  $\delta(G) + 1$  spanned on the vertices of  $U \cup \{z\}$ , and  $G_2$  is a complete graph of order  $n - \delta(G)$  spanned on  $V \cup \{z\}$ . Then  $z$  is a cutvertex, contradicting the assumption that  $G$  be 2-connected.

*Proof of part (ii).* Suppose there exists a 2-connected graph of odd order  $n \geq 3$ , in which  $d(x) + d(y) \geq n - 1$  for every pair of non-adjacent vertices  $x$  and  $y$ , that does not contain neither a Hamilton cycle nor a cycle of length  $n - 1$ . Let  $G$  be maximal such a graph of order  $n$ . By maximality of  $G$ ,  $G + xy$  contains a cycle of length at least  $n - 1$  whenever  $xy \notin E(G)$ . Hence  $G$  contains a path of length at least  $n - 2$  between any two of its non-adjacent vertices.

Pick a pair of non-adjacent vertices  $x$  and  $y$ . By a theorem of Pósa,  $G$  contains a Hamilton  $x - y$  path  $P$ , and hence, by Theorem 2.4, the sum  $d(x) + d(y)$  actually equals  $n - 1$ . Write  $P = u_1u_2 \dots u_n$ , where  $u_1 = x$  and  $u_n = y$ .

Put  $I_x = \{i : xu_{i+1} \in E(G), 1 \leq i \leq n - 1\}$  and  $I_y = \{i : u_iy \in E(G), 1 \leq i \leq n - 1\}$ . If  $I_x \cap I_y \neq \emptyset$ , say  $i_0 \in I_x \cap I_y$ , then  $G$  contains a Hamilton cycle

$$u_1u_{i_0+1}u_{i_0+2} \dots u_nu_{i_0}u_{i_0-1} \dots u_2u_1.$$

We may thus assume that  $I_x \cap I_y = \emptyset$ . Then, for every  $1 \leq i \leq n-1$ , either  $u_i$  is adjacent to  $y$  or else  $u_{i+1}$  is adjacent to  $x$ , because  $|I_x| + |I_y| = d(x) + d(y) = n-1$ . Let  $d = d(y)$  and let  $v_1, \dots, v_d = y$  be the vertices that lie on  $P$  next to the (respective) neighbours of  $y$ .

If there exists  $j < d$  such that  $v_j \notin N(y)$ , then  $v_j = u_{i_0}$  for some  $i_0 \in I_x$ . It follows that  $u_{i_0+1}$  is adjacent to  $x$ , and  $G$  contains a cycle of length  $n-1$  of the form

$$u_1 u_{i_0+1} u_{i_0+2} \dots u_n u_{i_0-1} u_{i_0-2} \dots u_2 u_1.$$

Therefore we can assume that

$$(\dagger) \quad v_1, \dots, v_{d-1} \text{ are all adjacent to } y.$$

Let  $z$  denote the furthestmost neighbour of  $y$  on  $P$ . It follows from  $(\dagger)$  that all the vertices between  $z$  and  $y$  on  $P$  are adjacent to  $y$ , and hence  $z = u_{n-d}$ .

Suppose  $N(v_j) \subset \{z, v_1, \dots, v_d\}$  for  $j \leq d$ . Then  $N(u_i) \subset \{u_1, \dots, u_{n-d-1}, z\}$  for  $i \leq n-d-1$ . Consequently,  $d(u_i) \leq n-d-1$ ,  $d(v_j) \leq d$ , and  $u_i v_j \notin E(G)$  for  $i \leq n-d-1$  and  $j \leq d$ . But then  $d(u_i) + d(v_j) \geq n-1$  yields

$$d(u_i) = n-d-1 \quad \text{and} \quad d(v_j) = d \quad \text{for } i = 1, \dots, n-d-1, j = 1, \dots, d.$$

Therefore, as in the proof of part (i), we get that  $G = G_1 \cup G_2$ , where  $G_1$  is a complete graph of order  $n-d$  spanned on the vertices  $\{u_1, \dots, u_{n-d-1}, z\}$  and  $G_2$  is a complete graph of order  $d+1$  on  $\{z, v_1, \dots, v_d\}$ . Then  $z$  is a cutvertex contradicting our assumptions on  $G$ .

It remains to consider the case of some  $v_{j_0}$  being adjacent to  $u_{i_0}$ , where  $i_0 \leq n-d-1$ . But then again  $G$  contains a Hamilton cycle

$$u_1 \dots u_{i_0} v_{j_0} \dots v_d v_{j_0-1} \dots u_{i_0+1} u_1.$$

For the proof of the last assertion of Theorem 1.4, suppose that  $n = 2k+1$  is odd and  $G$  is a non-hamiltonian 2-connected graph on  $n$  vertices, satisfying  $d(x) + d(y) \geq n-1$  for every pair of non-adjacent  $x$  and  $y$ . Then the  $n$ -closure of  $G$ ,  $G^* = Cl_n(G)$  is not hamiltonian either, by Theorem 2.5, and we have equality

$$d_{G^*}(x) + d_{G^*}(y) = n-1 \quad \text{whenever } xy \notin E(G^*).$$

Now, if  $\delta(G^*) < k = \frac{n-1}{2}$ , then  $n-1 - \delta(G^*) > \delta(G^*)$  and one can repeat the proof of part (i) to show that  $G^*$  contains a Hamilton cycle, which contradicts the assumptions on  $G$ .

Thus  $\delta(G^*) = \frac{n-1}{2}$ . Moreover,  $d_{G^*}(x) + d_{G^*}(y) = n-1 = 2k$  for  $xy \notin E(G^*)$  implies that  $d_{G^*}(x) = k$  or  $d_{G^*}(x) = n-1$  for every vertex  $x$ .

Suppose  $G^*$  is not maximal among the non-hamiltonian 2-connected graphs on  $n$  vertices. Then  $G^*$  has a pair of non-adjacent vertices  $x$  and  $y$  such that  $G^* + xy$  is contained in a maximal non-hamiltonian graph  $H$ . By maximality of  $H$ ,  $H + uv$  contains a Hamilton cycle for every  $uv \notin E(H)$ , so Theorem 2.4 implies that  $d_H(u) + d_H(v) = n-1$  for every  $uv \notin E(H)$ .

Notice that  $d_{G^*}(x) = k$ , as  $d_{G^*}(x) < n-1$ . Then  $d_H(x) \geq k+1$  and hence, for every  $v$  non-adjacent to  $x$  in  $G^*$ ,  $d_H(x) + d_H(v) \geq d_{G^*}(x) + 1 + d_{G^*}(v) > n-1$ , implying  $xv \in E(H)$ . Therefore  $H$  is obtained from  $G$  by increasing degrees of at least  $x$  and all its non-neighbours in  $G^*$ , that is, at least  $1 + (n-1-k) = k+1$  vertices. But then  $H$  contains at least  $k+1$  vertices of degree  $n-1$ , which

means that  $\delta(H) \geq k + 1 = \frac{n+1}{2}$ , and hence  $H$  is hamiltonian by Theorem 2.3; a contradiction.  $\square$

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