

Filter Products and Elementary Models of Homotopy Type Theory

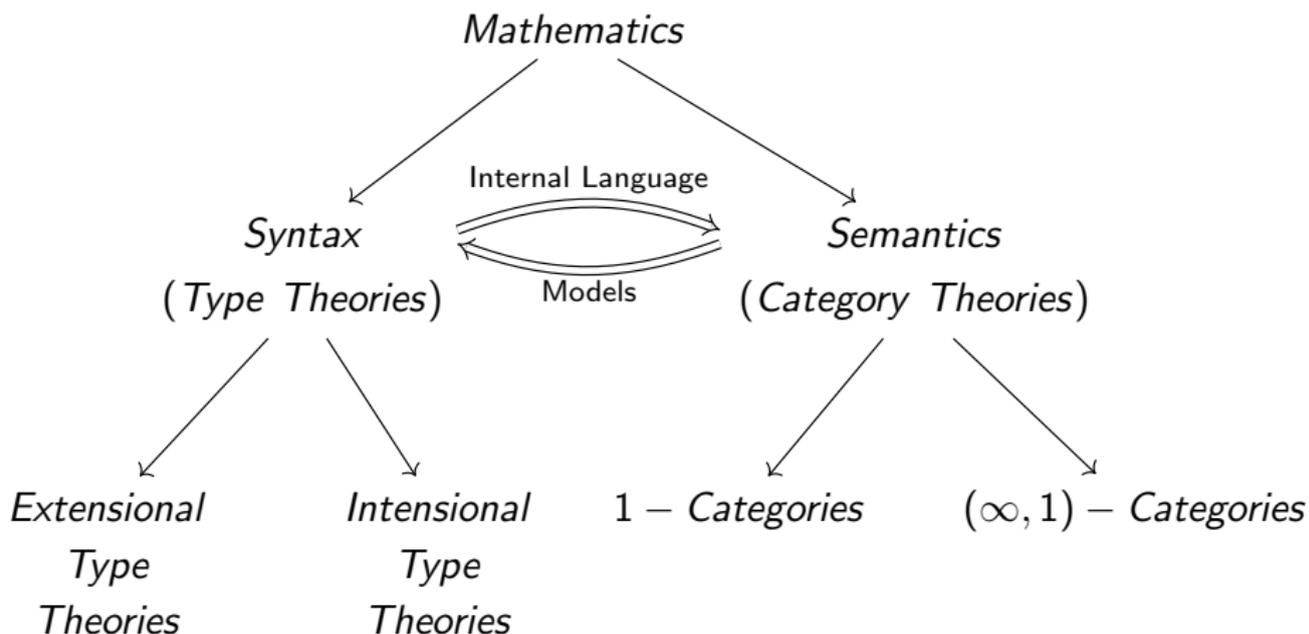
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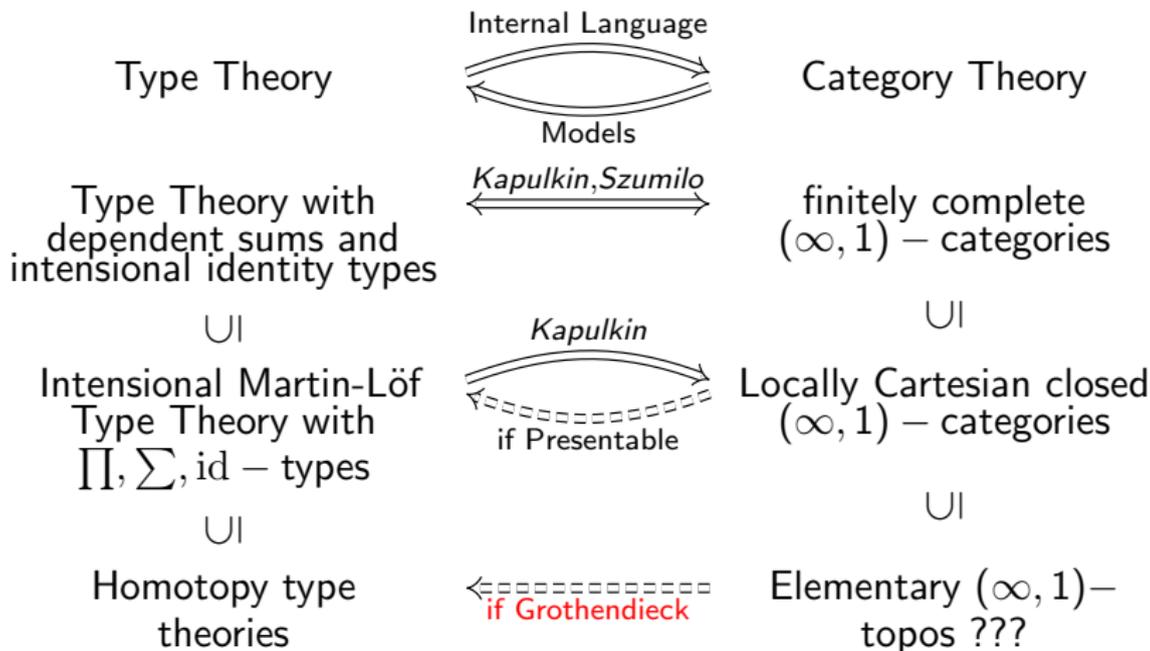


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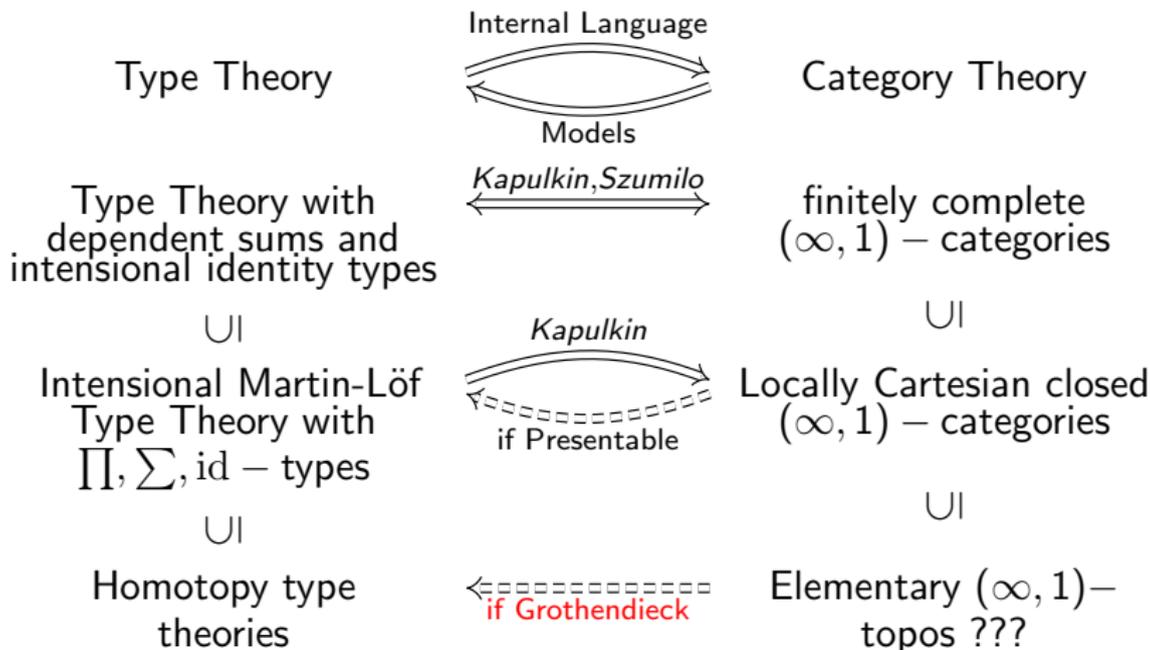
Syntax vs. Semantics



Intensional Type Theories vs. $(\infty, 1)$ -Categories



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Let's focus on the last line!

What's a model for HoTT: Vague Edition

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- 2 The type of natural numbers \rightsquigarrow natural number object.

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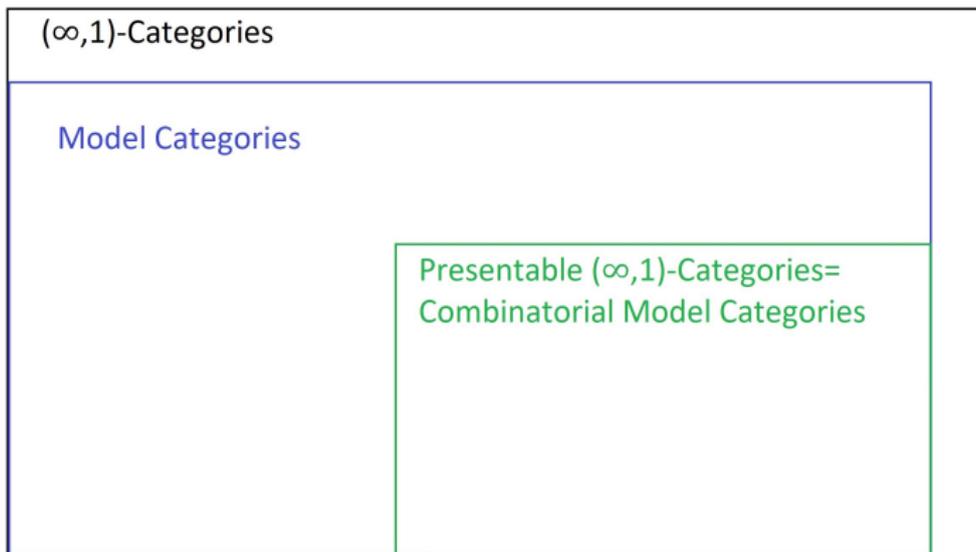
In some cases it's not really clear what that means:

- 3 Univalent universe?

Key problem: General $(\infty, 1)$ -categories are very non-strict!

Model Categories

The fix is an appropriate use of model categories!



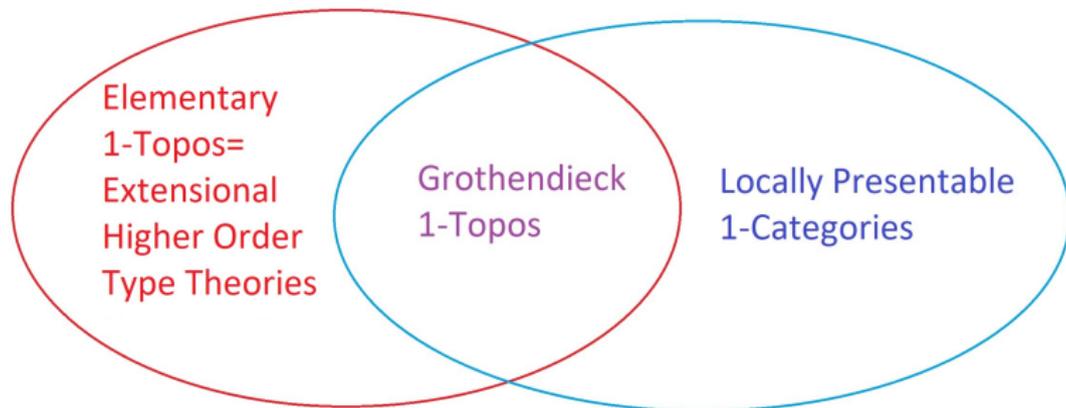
Model categories are strict 1-categories!

Axiomatizations of Model Categories

- 1 **Arendt, Kapulkin:** Introduce *logical model categories* and prove they model Σ , Π , id-types.
- 2 **Shulman:** Introduces *type-theoretic model categories* and prove their Π -types satisfy function extensionality.
- 3 **Shulman, Lumsdaine:** Introduce *good model categories*, which model certain higher inductive types (coproduct type, circle, ...), and *excellent model categories*, which model further higher inductive types (W -types, truncations, localizations, ...).
- 4 **Shulman:** Introduces *type-theoretic model toposes*, which is a model topos and is a special case of all the previous examples, but also models strict **univalent universes**.

Why Grothendieck Toposes?

Motivated by 1-Categories



Grothendieck Toposes and Model Toposes

Definition

A *Grothendieck 1-topos* is a category \mathcal{G} that fits into an adjunction

$$\text{Fun}(\mathcal{C}^{op}, \text{Set}) \begin{array}{c} \xrightarrow{a} \\ \perp \\ \xleftarrow{\quad} \end{array} \mathcal{G}$$

where \mathcal{C} **small**, a is **left-exact**.

if this fails \mathcal{G} is locally presentable

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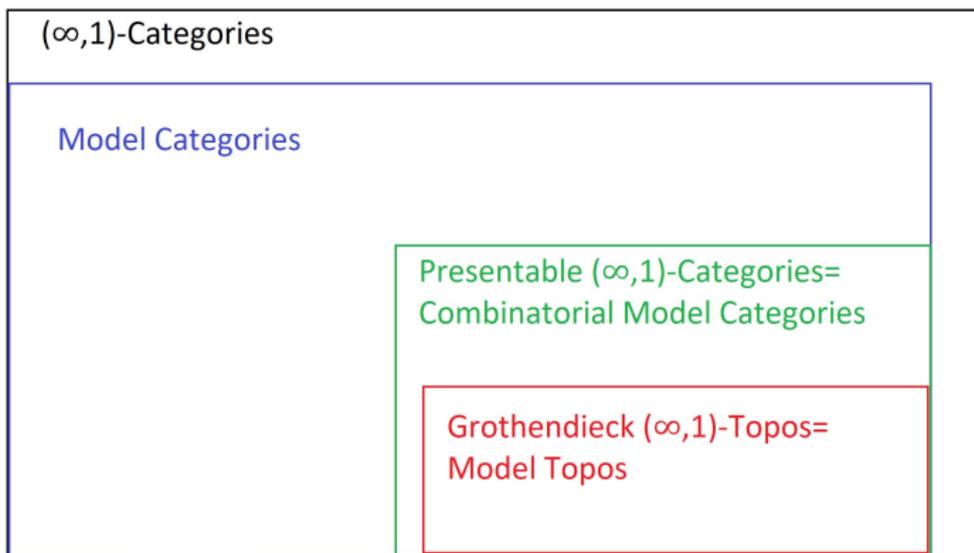
A *Grothendieck model topos* is a simplicial model category \mathcal{X} that fits into a Quillen adjunction:

$$\mathrm{Fun}(\mathcal{C}^{op}, \mathrm{sSet}^{Kan})^{proj} \begin{array}{c} \xrightarrow{a} \\ \leftarrow \perp \rightarrow \end{array} \mathcal{X}$$

where \mathcal{C} **small**, a is **left-exact**.

Model Topos vs. Grothendieck $(\infty, 1)$ -Topos

We generalize the previous diagram:



What is a Type-Theoretic Model Topos \mathcal{E} ?

Definition (Shulman)

- 1 Grothendieck 1-topos.
- 2 right proper, simplicial, combinatorial model structure with cofibrations monos.
- 3 simplicially locally Cartesian closed
- 4 locally representable, relatively acyclic notion of fibred structure that covers all fibrations

Note: It is in fact a **Grothendieck model topos**.

Univalent Universes in TTMT

Fix large enough κ .

$$\mathbb{Fib}^\kappa : \mathcal{E}^{op} \longrightarrow \mathbf{Grpd}$$

$$X \longmapsto ((\mathbb{Fib}/_X)^\kappa) \cong$$

underlying
groupoid

The last condition implies the existence of a cofibrant object \mathcal{U} and an *acyclic fibration*

$$\mathcal{E}(-, \mathcal{U}) \twoheadrightarrow \mathbb{Fib}^\kappa$$

Moreover, \mathcal{U} is fibrant and univalent.

Model Topos = Type-Theoretic Model Topos

Type-theoretic model toposes in fact recover all model toposes.

$$\text{Fun}(\mathcal{C}, \text{sSet}^{\text{Kan}}) \begin{array}{c} \text{proj} \\ \xrightarrow{a} \\ \xleftarrow{\perp} \\ \xrightarrow{\quad} \end{array} \mathcal{X}$$

Three steps:

- 1 The **Kan model structure** is a TTMT (also observed by Kapulkin-Lumsdaine).
- 2 **Injective** model structure on diagrams into TTMT is a TTMT.
- 3 **Left-exact Bousfield localizations** of a TTMT is a TTMT.

What about the non-presentable case?

We expect **non-presentable models** for homotopy type theories. However, arbitrary non-presentable $(\infty, 1)$ -categories don't come from model categories.

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We expect **non-presentable models** for homotopy type theories. However, arbitrary non-presentable $(\infty, 1)$ -categories don't come from model categories.

- 1 *General Case*: Embed every category in its presheaf category.
- 2 *Specific Case*: Use specific constructions.
 - **Realizability topos**
 - **Filter Product**

Elementary 1-Toposes

In order to move from the presentable to the non-presentable world we need to generalize our toposes.

Definition

An elementary 1-topos is a locally Cartesian closed category with subobject classifier.

Proposition

A locally presentable category is an elementary 1-topos if and only if it is a Grothendieck 1-topos.

Slogan: The “correct” generalization of Grothendieck 1-toposes.

Filters

Let's start with filters:

Definition

Let I be a set. A filter $\Phi \subseteq P(I)$ is a subset of the **power set** satisfying:

- 1 **Non-Empty:** $I \in \Phi$.
- 2 **Intersection Closed:** $J_1, J_2 \in \Phi \rightarrow J_1 \cap J_2 \in \Phi$.
- 3 **Upwards Closed:** $J_1 \in \Phi, J_1 \subseteq J_2 \rightarrow J_2 \in \Phi$.

We can think of elements in Φ as “large” subsets of I .

$$I \in \Phi \Leftrightarrow I^c \text{ "small"}$$

Filter Products

Definition

For a category \mathcal{C} define the **filter-product** $\prod_{\Phi} \mathcal{C}$ as:

- Obj: $(c_i)_{i \in I}$, $c_i \in \text{Obj}(\mathcal{C})$.
- Mor:

$$\text{Hom}_{\prod_{\Phi} \mathcal{C}}((c_i)_{i \in I}, (c'_i)_{i \in I}) = \left(\prod_{J \in \Phi} \prod_{i \in J} \text{Hom}_{\mathcal{C}}(c_i, c'_i) \right) / \sim$$

$(J, f_i: c_i \rightarrow c'_i)_{i \in J}$

$$(f_i)_{i \in J_1} \sim (g_i)_{i \in J_2} \Leftrightarrow \exists J_3 \subseteq J_1 \cap J_2, J_3 \in \Phi, \{i \in J_3 : f_i = g_i\} \in \Phi$$

So, morphisms that agree on a “large” index set are identified.

Filter Products in Elementary Topos Theory

Filter products are relevant in topos theory.

Theorem (Adelman-Johnstone 1982)

Let \mathcal{E} be an elementary 1-topos, I a set and Φ a filter. Then $\prod_{\Phi} \mathcal{E}$ is also an elementary 1-topos.

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Example

Let $\mathcal{E} = \mathbf{Set}$, $I = \mathbb{N}$ and Φ the **Fréchet filter** (of cofinite sets). Then $\prod_{\Phi} \mathbf{Set}$ is a non-Grothendieck elementary topos.

finite = J "small" iff $J^c \in \Phi$

Generalize to Model Categories

We want to generalize the example from categories to model categories!

We need to generalize the definition appropriately:

Type-Theoretic Elementary Model Topos

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	TT Grothendieck MT	TT Elementary MT
(1)	Grothendieck topos	
(2)	right proper, simplicial, combinatorial , cofibrations are monos	
(3)	simplicially lcc	
(4)	notion of fibred structure \mathbb{F} locally representable relatively acyclic ? $ \mathbb{F} = \mathbb{F}ib$ ✓	

Type-Theoretic Elementary Model Topos

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(1)	Grothendieck topos	elementary topos
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(3)	simplicially lcc	simplicially lcc
(4)	notion of fibred structure \mathbb{F} locally representable relatively acyclic $ \mathbb{F} = \mathbb{F}ib$	$\mathbb{F}ib$ has a fibrant-cofibrant univalent universe

What does it mean: “has a universe”?

$$\text{Fib}^\kappa: \mathcal{E}^{\text{op}} \longrightarrow \text{Grp} \\ X \longmapsto (\text{Fib}/X) \simeq$$

\mathcal{E} has a universe, if there is a filtration

$$\mathcal{E}^{\kappa_1} \subseteq \mathcal{E}^{\kappa_2} \subseteq \dots$$

of \mathcal{E} such that for all \mathcal{E}^κ in the filtration there exists a fibrant-cofibrant, univalent universe \mathcal{U} , meaning a acyclic ~~trivial~~ fibration:

$$\mathcal{E}(-, \mathcal{U}) \twoheadrightarrow \text{Fib}^\kappa.$$

Is this the best we can do?

I don't know!

- 1 It **generalizes** type-theoretic Grothendieck model topos.
- 2 It still **includes** many relevant examples:
 - Logical model categories
 - Type-theoretic model categories
 - good model categories
 - but **not** excellent model categories

In particular, it interprets **Martin-Löf Type theory** with Σ -types, Π -types with function extensionality, identity types, the natural numbers type, the sphere types S^n , universe types, ...

- 3 It has **non-trivial examples**.

Filter-Product Model Structure

Theorem (R)

Let \mathcal{M} be a model structure (with **finite** (co)limits), I a set and Φ a filter on I . Then there is a model structure on $\prod_{\Phi} \mathcal{M}$ given by

$$(f_i)_{i \in J} \in \mathcal{F} \Leftrightarrow \{i \in J : f_i \in \mathcal{F}\} \in \Phi$$

where \mathcal{F} is one of the classes of fibrations/cofibration/weak equivalences.

The proof is routine checking.

Remark

Fun fact: We really need all three conditions of a filter!

Transfer of Properties

The following properties will transfer from \mathcal{M} to $\prod_{\Phi} \mathcal{M}$:

- finite (co)limits
- Cartesian closure
- left/right proper
- simplicial
- compatibility with Cartesian closure
- cofibrations monomorphism

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The following will **not** transfer:

- infinite (co)limits
- local presentability
- cofibrantly generated

Comparison with $(\infty, 1)$ -Version

We also have following comparison theorem.

Theorem (R)

Let \mathcal{M} be a simplicial model category, I a set and Φ a filter. Then we have an equivalence of $(\infty, 1)$ -categories

$$N\left(\prod_{\Phi} \mathcal{M}\right) \simeq \prod_{\Phi} N(\mathcal{M})$$

Constructing Elementary Models of HoTT

Theorem (R)

Let \mathcal{E} be a type-theoretic elementary model topos, I a set and Φ a filter. Then $\prod_{\Phi} \mathcal{E}$ is a type-theoretic elementary model topos.

This directly generalizes the result by Adelman-Johnstone.

Example I

We are now finally in a position to put the theory in practice and give examples.

Example

Let

- $\mathcal{E} = \text{sSet}$ with $\mathcal{K}\text{an}$ model structure
- $I = \mathbb{N}$
- $\Phi = \text{Fréchet filter}$ (cofinite subsets)

$I \in \Phi \Leftrightarrow I^c$ finite

Then, by the previous theorem, $\prod_{\Phi} \mathcal{K}\text{an}$ is a type-theoretic elementary model topos.

Claim: It's not Grothendieck!

Example II

- ① The underlying category $\prod_{\phi} \mathbf{sSet}$ is not locally presentable. So, in particular, $\prod_{\phi} \mathbf{Kan}$ is not combinatorial.
- ② It does not even have infinite colimits.
- ③ The natural number object is non-standard.

$$\frac{11}{\pi} \quad 1 \quad \mathbb{N} \times \mathbb{N} \quad 0$$

$$\mathbf{Kan} \times \mathbf{Kan} \quad (n, m) \in \mathbb{N} \times \mathbb{N}$$

$$n \times \underline{(1, 0)} + m \times (0, 1)$$

Where do we go from here?

- These results don't generalize to other elementary models!
- Could be an application of HoTT to non-classical algebraic topology
 - HoTT indexes homotopy groups by the internal natural number object.
 - HoTT proofs of algebraic topological results still hold
 - ...

Proof I

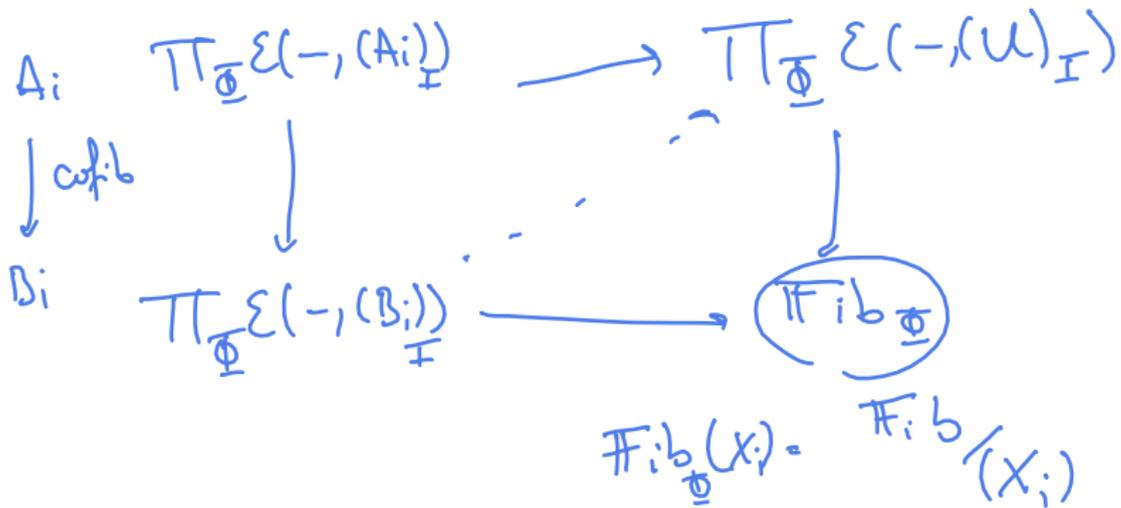
Need to check four conditions:

- ✓ ① $\prod_{\phi} \mathcal{E}$ is an elementary 1-topos.
- ✓ ② $\prod_{\phi} \mathcal{E}$ has a right proper, simplicial model structure where cofibrations are monos
- ✓ ③ $\prod_{\phi} \mathcal{E}$ is simplicially locally Cartesian closed.
- ④ \mathbf{Fib}_{ϕ} has a fibrant-cofibrant univalent universe

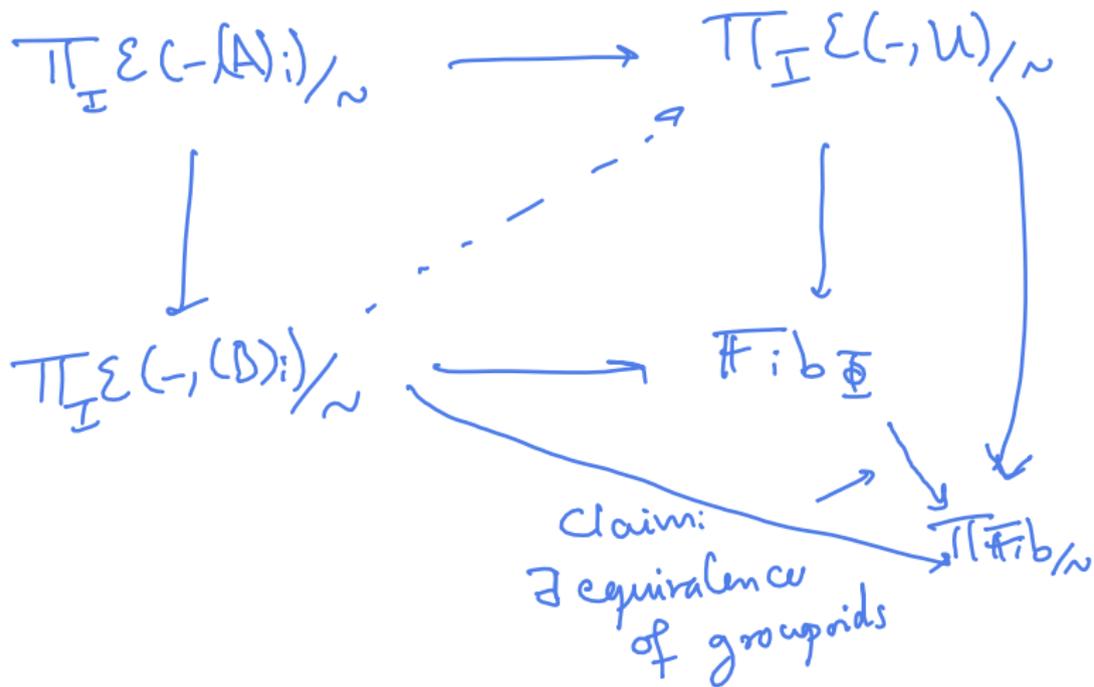
The first three follow from previous analysis.

Proof II

Claim: If \mathcal{U} is a universe in \mathcal{E} , then $(\mathcal{U})_{\mathcal{I}}$ is also a universe in $\Pi_{\mathcal{I}}\mathcal{E}$



Proof III



The End!

Thank you!

Questions?