

Path spaces of pushouts via a zigzag construction

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Plan for the talk

- ▶ Discuss the background of the problem of understanding path spaces of pushouts
- ▶ Present the zigzag construction. First in a ‘type-theoretic’ style, then in a ‘diagrammatic’ style due to Christian Sattler.
- ▶ Discuss ‘convergence behaviour’ and applications to truncatedness of pushouts, Blakers–Massey, combinatorial group theory

Pushouts

We work informally in HoTT with a univalent universe U closed under *pushouts*.

Given

- ▶ $A : U$
- ▶ $B : U$
- ▶ $R : A \rightarrow B \rightarrow U$

the pushout $A +_R B : U$ is freely generated by

- ▶ $\text{inl} : A \rightarrow A +_R B$
- ▶ $\text{inr} : B \rightarrow A +_R B$
- ▶ $\text{glue} : (a : A) (b : B) \rightarrow R a b \rightarrow \text{inl } a = \text{inr } b.$

The problem

We know that any element of $A +_R B$ is merely either of the form $\text{inl } a$ or $\text{inr } b$.

But what are the identity types / path spaces?

- ▶ $\text{inl } a_0 = \text{inr } b$?
- ▶ $\text{inl } a_0 = \text{inl } a$?
- ▶ $(\text{inr } b = \text{inr } b')$?

For most type formers (Σ -, Π -, W -, M -types, univalent universes, n -truncations, sequential colimits), identity types are easy to describe.

Identity types of pushouts are 'complicated,' like $\Omega^m S^n$.

Some answers

Aside from the generating paths $R a b \rightarrow \text{inl } a = \text{inr } b$ and their inverses one also has anything built up from a zigzag of generating paths.

We have to quotient by inverse laws $(\text{glue } r) \cdot (\text{glue } r)^{-1} = \text{refl}$, but how exactly?

Can be made precise if we are interested in set-truncation $\|\text{inl } a = \text{inr } b\|_0$.¹

¹Favonia, Shulman: The Seifert-van Kampen Theorem in Homotopy Type Theory; 2016

Some answers

Kraus and von Raumer² gave a universal property for path spaces of pushouts:

For any type $X : U$ with $x : X$, the type family $X \rightarrow U, y \mapsto x = y$ is freely generated by $\text{refl} : x = x$.

A type family $A +_R B \rightarrow U$ is the same thing as the data

- ▶ $P : A \rightarrow U$
- ▶ $Q : B \rightarrow U$
- ▶ $e : (a : A) (b : B) \rightarrow R a b \rightarrow P a \simeq Q b$.

Thus the triple (P, Q, e) freely generated by a point $p : P a_0$ has $P a \simeq (\text{inl } a_0 = \text{inl } a)$, $Q b \simeq (\text{inl } a_0 = \text{inr } b)$, and e corresponding to post-composition by glue .

²Path Spaces of Higher Inductive Types in Homotopy Type Theory; 2019

Some answers

This is a nice description and easy to prove, but not immediately useful for many purposes.

A priori describes path spaces of non-recursive HITs (pushouts) as recursive HITs (much more complex).

There is a close analogy with the James construction.³

³Brunerie: The James Construction and $\pi_4(\mathbb{S}^3)$ in Homotopy Type Theory; 2018

James construction

If $X : \mathcal{U}$, $x_0 : X$ is a *pointed connected* type, writing $\Sigma X := 1 +_X 1$, we have $JX := \Omega \Sigma X$ freely generated by

- ▶ a term $\varepsilon : JX$
- ▶ a map $\alpha : X \rightarrow JX \rightarrow JX$
- ▶ with $\delta : (j : JX) \rightarrow \alpha(x_0, j) = j$.

Also a recursive description, but can be ‘unrecursified’, so $JX = \operatorname{colim}_{n \rightarrow \infty} J_n X$.

α unrecursifies to $\alpha_n : X \rightarrow J_n X \rightarrow J_{n+1} X$ with a naturality condition.

Dealing with equivalences

We would like to follow to the same strategy to ‘unrecursify’ Kraus–von Raumer’s description of path spaces of pushouts, so $P a = \operatorname{colim}_{n \rightarrow \infty} P_n a$ and $Q b = \operatorname{colim}_{n \rightarrow \infty} Q_n b$.

But how to unrecursify $P(a) \simeq Q(b)$ for $r : R a b$?

Several ways to express equivalence constructors in HITs:

Kraus and von Raumer use biinvertible maps.

Rijke, Shulman, and Spitters⁴ use *path-split* maps.

One could probably unrecursify these, but they seem to give the *wrong* sequence P_n .

We follow a different route.

⁴Modalities in homotopy type theory; 2020

Interleaving sequences

Lemma

Given sequences $P_0 \rightarrow P_2 \rightarrow \dots$ and $Q_1 \rightarrow Q_3 \rightarrow \dots$ and a commutative diagram

$$\begin{array}{ccccccc} P_0 & \longrightarrow & P_2 & \longrightarrow & \dots & & \\ & \searrow & \nearrow & \searrow & \nearrow & & \\ & & Q_1 & \longrightarrow & Q_3 & \longrightarrow & \dots \end{array}$$

we have an equivalence $\operatorname{colim}_{n \rightarrow \infty} P_n \simeq \operatorname{colim}_{n \rightarrow \infty} Q_n$.

Proof sketch.

Consider the sequence $P_0 \rightarrow Q_1 \rightarrow P_2 \rightarrow \dots$ and its colimit.

Omitting every other term, we get the two sequences

$P_0 \rightarrow P_2 \rightarrow \dots$ and $Q_1 \rightarrow Q_3 \rightarrow \dots$ we started with.

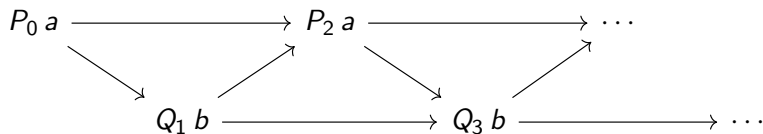
Thus all three sequences have the same colimit. □

The zigzag construction

Let $A, B : U$, $R : A \rightarrow B \rightarrow U$, $a_0 : A$ as before.

We define $P_0, P_2, \dots : A \rightarrow U$ and $Q_1, Q_3, \dots : B \rightarrow U$ freely so that we have

- ▶ a map $P_n a \rightarrow P_{n+2} a$ for $a : A$, n even,
- ▶ a map $Q_n b \rightarrow Q_{n+2} b$ for $b : B$, n odd
- ▶ a term of $P_0(a_0)$,
- ▶ for $a : A$, $b : B$, $r : R a b$, an interleaving diagram.



The zigzag construction

More concretely, we have

- ▶ $P_0 a := (a_0 = a)$
- ▶ $Q_1 b := R a_0 b$
- ▶ $P_{n+2} a$ given by a pushout square

$$\begin{array}{ccc} (b : B) \times R a b \times P_n a & \longrightarrow & (b : B) \times R a b \times Q_{n+1} b \\ \downarrow & & \downarrow \\ P_n a & \xrightarrow{\quad \Gamma \quad} & P_{n+2} a \end{array}$$

- ▶ $Q_{n+2} b$ is given by the analogous pushout square

$$\begin{array}{ccc} (a : A) \times R a b \times Q_n b & \longrightarrow & (a : A) \times R a b \times P_{n+1} a \\ \downarrow & & \downarrow \\ Q_n b & \xrightarrow{\quad \Gamma \quad} & Q_{n+2} b \end{array}$$

The zigzag construction

Theorem

With notation as before,

$(\text{inl } a_0 = \text{inl } a) \simeq \text{colim}_{n \rightarrow \infty} P_n a$ for $a : A$ and

$(\text{inl } a_0 = \text{inr } b) \simeq \text{colim}_{n \rightarrow \infty} Q_n b$ for $b : B$.

Proof sketch.

Write $P_\infty a := \text{colim}_{n \rightarrow \infty} P_n a$.

Then $e r : P_\infty a \simeq Q_\infty b$ for $r : R a b$ by the interleaving diagram.

Now (P_∞, Q_∞, e) is freely generated by a term of $P_\infty a_0$ essentially by construction.

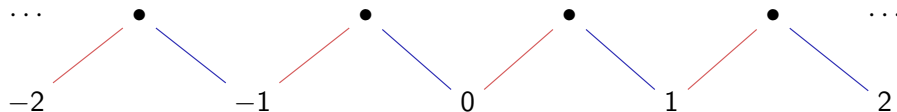
So we can appeal to Kraus–von Raumer’s characterisation of the path spaces. □

An example: ΩS^1

We have $S^1 \simeq 1 +_2 1$ where $2 = \{B, R\}$.⁵

Writing $N = \text{inl} \star$ and $S = \text{inr} \star$, the construction describes $N = N$ and $N = S$ as sequential colimits. How?

We picture $N = N$ as the bottom row above and $N = S$ as the top row. The filtrations (i.e. types P_n and Q_n) are given by intervals centred on 0.



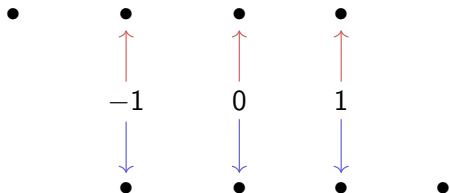
⁵Note that a relation $R : 1 \rightarrow 1 \rightarrow U$ is just a type.

An example: ΩS^1

The following pushout diagram describes P_4 :

$$\begin{array}{ccc} 2 \times P_2 & \longrightarrow & 2 \times Q_3 \\ \downarrow & & \downarrow \\ P_2 & \longrightarrow & P_4 \end{array}$$

This expresses that $\{-2, -1, 0, 1, 2\}$ is given by gluing two four-element sets along $\{-1, 0, 1\}$.



A diagrammatic perspective

An alternative presentation of the construction due to Christian Sattler avoids type theory-style indexing.

Say given a span of spaces $A \leftarrow R \rightarrow B$.

Given a map $Y \rightarrow A +_R B$ (e.g. $\text{inl} : A \rightarrow A +_R B$) we seek to understand the pullback of $A \leftarrow R \rightarrow B$ along $Y \rightarrow A +_R B$.

In general suppose we have a span $P_0 \leftarrow T_0 \rightarrow Q_0$ over the first one. We describe the pullback of the first span along the induced map $P_0 +_{T_0} Q_0 \rightarrow A +_R B$.

To this end we construct a sequence of overspans

$(P_n \leftarrow T_n \rightarrow Q_n)_{n:\mathbb{N}}$ and take the colimit $P_\infty \leftarrow T_\infty \rightarrow Q_\infty$.

A diagrammatic perspective

Descent for pushouts means that if both squares below are pullback squares then the top span is the pullback of the bottom one along $P_\infty +_{T_\infty} Q_\infty \rightarrow A +_R B$.

$$\begin{array}{ccccc} P_\infty & \longleftarrow & T_\infty & \longrightarrow & Q_\infty \\ \downarrow & & \lrcorner & & \downarrow \\ A & \longleftarrow & R & \longrightarrow & B \end{array}$$

We can in turn ensure that these squares are pullback squares – in short that $P_\infty \leftarrow T_\infty$ and $T_\infty \rightarrow Q_\infty$ are cartesian – by ensuring that $P_n \leftarrow T_n$ and $T_n \rightarrow Q_n$ are each cartesian for infinitely many n .

(This uses commutativity of pullbacks and sequential colimits.)

A diagrammatic perspective

Given $P \leftarrow T \rightarrow Q$ a span over $A \leftarrow R \rightarrow B$ we construct a span $P' \leftarrow T' \rightarrow Q'$ in between the above two such that

- ▶ $P \rightarrow P'$ is an equivalence,
- ▶ $P' \leftarrow T'$ is cartesian over $A \leftarrow R$,
- ▶ $P +_T Q \rightarrow P' +_{T'} Q'$ is an equivalence,

as follows.

$$\begin{array}{ccccc} P & \longleftarrow & T & \longrightarrow & Q \\ \downarrow \text{id} & & \downarrow & \lrcorner & \downarrow \\ P' & \longleftarrow & T' & \longrightarrow & Q' \\ \downarrow & & \downarrow & \lrcorner & \downarrow \\ A & \longleftarrow & R & \longrightarrow & B \end{array}$$

A diagrammatic perspective

Now starting from $P_0 \leftarrow T_0 \rightarrow Q_0$ one can iterate the previous construction, alternating between making $P_n \leftarrow T_n$ cartesian and making $T_n \rightarrow Q_n$ cartesian.

The pushout is unchanged in each step so is unchanged also in the colimit as $n \rightarrow \infty$.

Thus $P_\infty \leftarrow T_\infty \rightarrow Q_\infty$ is precisely the pullback of $A \leftarrow R \rightarrow B$ along $P_0 +_{T_0} Q_0 \rightarrow A +_R B$.

We have e.g. $A = A +_0 0$, and $1 = 1 +_0 0$.

The first few steps

$$\begin{array}{ccccc}
 A & \longleftarrow & 0 & \xrightarrow{\text{cart}} & 0 \\
 \downarrow & & \downarrow & & \downarrow \\
 A & \longleftarrow & R & \xrightarrow{\quad \sqsupset \quad} & R \\
 \downarrow & & \downarrow & & \downarrow \\
 A +_R R \times_B R & \longleftarrow & R \times_B R & \xrightarrow{\text{cart}} & R \\
 \downarrow & & \downarrow & & \downarrow \\
 A +_R R \times_B R & \longleftarrow & (A +_R R \times_B R) \times_A R & \xrightarrow{\quad \sqsupset \quad} & \cdots \\
 \downarrow & & \downarrow & & \downarrow \\
 \vdots & & \vdots & & \vdots
 \end{array}$$

Convergence behaviour

It will be fruitful to analyse the convergence behaviour of the construction.

How well does P_n approximate P_∞ ?

How 'far' is the map $P_n \rightarrow P_{n+2}$ from being an equivalence?

We are particularly interested in $P_0 a \rightarrow P_\infty a$, corresponding to $\text{ap}_{\text{inl}} : (a_0 = a) \rightarrow (\text{inl } a_0 = \text{inl } a)$,
and in $Q_1 b \rightarrow Q_\infty b$, corresponding to
 $\text{glue} : R a_0 b \rightarrow (\text{inl } a_0 = \text{inr } b)$.

Informal explanation of convergence behaviour

The fibres of the map $P_n \rightarrow P_{n+2}$ express when a zigzag of length at most $n + 2$ has length at most n .

This happens when one can *reduce* an pair of adjacent edges.

This is controlled by the identity types of $(a : A) \times R a b$ for $b : B$ and of $(b : B) \times R a b$ for $a : A$,

or equivalently by the diagonals of $R \rightarrow B$ and $R \rightarrow A$.

It is enough to reduce *some* adjacent pair in the zigzag.

This is why *joins* of these identity types show up.

Formal analysis of convergence behaviour

We want to understand the map $P_n a \rightarrow P_{n+2} a$.

It is defined as a pushout of

$$(b : B) \times R a b \times P_n a \rightarrow (b : B) \times R a b \times Q_{n+1} b.$$

So suffices to understand $P_n a \rightarrow Q_{n+1} b$ given $r : R a b$.

Theorem

*For $r : R a b$ and n even, the map $P_n a \rightarrow Q_{n+1} b$ is a pushout of a map f such that all fibres of f are of the form $X * Y$ where X is the fibre of a map $Q_{n-1} b' \rightarrow P_n a$ given by $r' : R a b'$ and Y is $(b, r) = (b', r')$.*

*Similarly, for $r : R a b$ and n odd, the map $Q_n b \rightarrow P_{n+1} a$ is a pushout of a map g such that all fibres of g are of the form $X * Y$ where X is the fibre of a map $P_{n-1} a' \rightarrow Q_n b$ given by $r' : R a' b$ and Y is $(a, r) = (a', r')$.*

The key theorem

Theorem

Let C_{-1}, C_0, C_1, \dots be classes of maps of types such that

- ▶ Each class is determined fibrewise: there is a class T_n of types such that C_n consists of all maps whose fibres are all in T_n .
- ▶ Each class is closed under pushouts.
- ▶ T_{-1} contains the empty type.
- ▶ For each $n \geq 0$, T_n contains any type that is a join of a type in T_{n-1} with an identity type in $(a : A) \times R a b$ if n is even, $b : B$ or an identity type in $(b : B) \times R a b$ if n is odd, $a : A$.

Then the maps $P_n a \rightarrow Q_{n+1} b$ and $P_n a \rightarrow P_{n+2} a$ lie in C_n for n even and $Q_n b \rightarrow P_{n+1} a$, $Q_n b \rightarrow Q_{n+2} b$ lies in C_n for n odd.

If moreover $C_{n+1} \subseteq C_n$ for all n and C_n is closed under transfinite composition, then the same holds for $P_n a \rightarrow P_\infty a$ and $Q_n b \rightarrow Q_\infty b$.

Pushouts of embeddings

Suppose $R \rightarrow B$ is an embedding.

Then we can take C_n to be the class of all equivalences for $n \geq 0$.

This shows that ap_{inl} is an equivalence i.e. $\text{inl} : A \rightarrow A +_R B$ is an embedding i.e. embeddings are closed under pushouts.

Also glue is an equivalence so the pushout square is a pullback square.

If $R \rightarrow A$ is an embedding we can take C_n to be all equivalences for $n \geq 1$ to see that $P_2 a \simeq (\text{inl } a_0 = \text{inl } a)$.

Pushouts of 0-truncated spans

Theorem

Suppose $R \rightarrow A$ and $R \rightarrow B$ are both 0-truncated i.e. their diagonals are embeddings. Then the same holds for inl and inr , and glue is an embedding.

Proof.

Take each C_n to consist of all embeddings. □

The same result holds if we replace 'embedding' with 'complemented' / 'decidable embedding' throughout.

Truncatedness of pushouts

Corollary

If $R \rightarrow A$ and $R \rightarrow B$ are both 0-truncated and A, B are both n -truncated with $n \geq 1$ then $A +_R B$ is also n -truncated.

Proof.

To be n -truncated means that Ω^{n+1} is contractible at each point, and inl, inr induce equivalences already on Ω^2 . □

So the suspension of a set, or any other pushout of sets, is 1-truncated. This resolves an open question from the HoTT book.

Some group theory

The following observation is due to Buchholtz, de Jong, and Rijke.

Theorem

Given a parallel pair of group embeddings $H \rightrightarrows G$, we have that G embeds in the associated HNN extension $G*_H$.

Proof.

The coequaliser of $BH \rightrightarrows BG$ is a delooping of $G*_H$. Equivalently this is a pushout:

$$\begin{array}{ccc} BH + BH & \longrightarrow & BH \\ \downarrow & \lrcorner & \downarrow \\ BG & \longrightarrow & B(G*_H) \end{array}$$

□

The proof is directly constructive and avoids combinatorial reasoning about words.

The Blakers–Massey theorem

Theorem

Let $k, l \geq 0$ be integers such that the diagonal of $R \rightarrow A$ is k -connected and the diagonal of $R \rightarrow B$ is l -connected. Then glue is $(k + l + 2)$ -connected.

Proof.

Take C_n to be the class of $((l + 2) + (k + 2) + (l + 2) + \dots - 2)$ -connected maps, where the sum contains $n + 1$ terms.

Then glue lies in C_1 which is the class of $(l + 2 + k + 2 - 2)$ -connected maps. □

This directly generalises a corresponding argument for the James construction.

A rough preprint with some more details is available online at
dwarn.se/po-paths.pdf

Thanks for listening!